

## $s$ -th power of Fibonacci number of the form

$$2^a + 3^b + 5^c$$

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**Received:** 25 February 2019

**Revised:** 6 August 2019

**Accepted:** 10 October 2019

**Abstract:** In this paper, we solve the Diophantine equation  $F_n^s = 2^a + 3^b + 5^c$ , where  $a, b, c$  and  $s$  are positive integers with  $1 \leq \max\{a, b\} \leq c$ .

**Keywords:** Fibonacci numbers, Linear forms in logarithms, Reduction method,  $s$ -th power.

**2010 Mathematics Subject Classification:** 11B39, 11J86.

## 1 Introduction

Let  $(F_n)_{n \geq 0}$  be the Fibonacci sequence given by the relation  $F_n = F_{n-1} + F_{n-2}$  with  $F_0 = 0$ ,  $F_1 = 1$  for all  $n \geq 2$ . It has many amazing combinatorial identities (see [7]). Put  $\alpha = (1 + \sqrt{5})/2$  and  $\beta = (1 - \sqrt{5})/2$ . Then the well-known Binet formula

$$F_n = \frac{\alpha^n - \beta^n}{\alpha - \beta} \quad (1)$$

holds for  $n \geq 0$ .

The problem of finding the different types of numbers among the terms of a linear recurrence has a long history. One of the popular results by Bageaud, Mignotte and Siksek [2] is that the integers 0, 1, 8, 144 among the Fibonacci numbers and the integers 1, 4 among the Lucas

numbers (an associated sequence of Fibonacci) can be written in the form  $y^t$  where  $t > 1$ . Szalay and Luca [4] showed that there are only finitely many quadruples  $(n, a, b, p)$  such that  $F_n = p^a \pm p^b + 1$  where  $p$  is a prime number. Marques and Togbé [5] determined the Fibonacci numbers and the Lucas numbers of the form  $2^a + 3^b + 5^c$  under  $1 \leq \max\{a, b\} \leq c$ . Bertók, Hajdu, Pink and Rábani [1] removed this condition. Namely, they gave full solutions of the equation

$$U_n = 2^a + 3^b + 5^c,$$

where  $U_n$  is the  $n$ -th Fibonacci, Lucas, Pell or Pell–Lucas number. We refer to the paper of Shorey and Stewart [9] for pure powers in recurrence sequences and some related Diophantine equations.

In this work, we generalize the problem of Marques and Togbé. We solve the Diophantine equation

$$F_n^s = 2^a + 3^b + 5^c \tag{2}$$

for  $s \geq 1$  integer and  $1 \leq \max\{a, b\} \leq c$ . Our result is following,

**Theorem 1.1.** *The solution of the equation (2) is  $(n, s, a, b, c) = (3, 5, 2, 1, 2)$ .*

## 2 Auxiliary results

Before going further, we present several lemmas. The following lemma was given by Matveev [6].

**Lemma 2.1.** *Let  $\mathbb{K}$  be a number field of degree  $D$  over  $\mathbb{Q}$ ,  $\gamma_1, \gamma_2, \dots, \gamma_t$  be positive real numbers of  $\mathbb{K}$ , and  $b_1, b_2, \dots, b_t$  be rational integers. Put*

$$B \geq \max\{|b_1|, |b_2|, \dots, |b_t|\},$$

and

$$\Lambda := \gamma_1^{b_1} \cdots \gamma_t^{b_t} - 1.$$

Let  $A_1, \dots, A_t$  be real numbers such that

$$A_i \geq \max\{Dh(\gamma_i), |\log \gamma_i|, 0.16\}, \quad i = 1, \dots, t.$$

Then, assuming that  $\Lambda \neq 0$ , we have

$$\begin{aligned} |\Lambda| &> \exp(-1.4 \times 30^{t+3} \times t^{4.5} \times D^2) \\ &\quad \times (1 + \log D)(1 + \log B) A_1 \dots A_t \end{aligned} \tag{3}$$

As usual, in the above lemma, the logarithmic height of the algebraic number  $\eta$  is defined as

$$h(\eta) = \frac{1}{d} \left( \log a_0 + \sum_{i=1}^d (\max\{|\eta^{(i)}|, 1\}) \right)$$

with  $d$  being the degree of  $\eta$  over  $\mathbb{Q}$  and  $(\eta^{(i)})_{1 \leq i \leq d}$  being the conjugates of  $\eta$  over  $\mathbb{Q}$ .

Application of the Matveev theorem gives the large upper bound. In order to reduce this bound, we use the following lemma.

**Lemma 2.2.** *Suppose that  $M$  is a positive integer. Let  $p/q$  be a convergent of the continued fraction expansion of the irrational number  $\gamma$  such that  $q > 6M$  and  $\epsilon = \|\mu q - M\| \|\gamma q\|$ , where  $\mu$  is a real number and  $\|\cdot\|$  denotes the distance from the nearest integer. If  $\epsilon > 0$ , then there is no solution to the inequality*

$$0 < m\gamma - n + \mu < AB^{-m}$$

*in positive integers  $m$  and  $n$  with*

$$\frac{\log(Aq/\epsilon)}{\log B} \leq m < M.$$

The following lemma is in the paper [3] (the case  $k = 2$ ).

**Lemma 2.3.** *For every positive integer  $n \geq 2$ , we have*

$$\alpha^{n-2} \leq F_n \leq \alpha^{n-1},$$

*where  $\alpha$  is the dominant root of the characteristic equation  $x^2 - x - 1 = 0$ .*

**Lemma 2.4.** *There is no solution of the equation*

$$2^s = 2^a + 3^b + 5^c \tag{4}$$

*for  $1 \leq \max\{a, b\} \leq c$ ,  $c \geq 6$  and  $s$  being positive integers.*

*Proof.* By (4) together with the facts that  $2 < \sqrt{5}$  and  $3 < 5^{0.7}$ , we get

$$|1 - 2^s 5^{-c}| < \frac{2}{(1.6)^c}. \tag{5}$$

We take  $\alpha_1 := 2$ ,  $\alpha_2 := 5$ ,  $b_1 := s$ ,  $b_2 := c$ . For this choice  $D = 1$ ,  $t = 2$ ,  $B = s$  and  $A_1 = 0.7 > \log 2$ ,  $A_2 = 1.61 > \log 5$ . The Lemma 2.1 yields that

$$\exp(C \cdot (1 + \log s)) < |1 - 2^s 5^{-c}| < \frac{2}{(1.6)^c}, \tag{6}$$

where  $C := 1.4 \cdot 30^5 \cdot 2^{4.5} \cdot 0.7 \cdot 1.61$ . Since  $2^s < 5^{c+1}$ , we have that  $0.4s < c+1$ . So, the inequality

$$s < 2.5 \cdot 10^{11}$$

is obtained. Let  $z := |s \log 2 - c \log 5|$ . Note that (5) can be written as

$$|1 - e^z| < \frac{3.2}{(1.2)^s},$$

since  $0.4s < c + 1$  holds. Since  $1 < 2^s 5^{-c}$ , then  $z > 0$  holds. We obtain that

$$0 < |s \log 2 - c \log 5| < |1 - e^z| < \frac{3.2}{(1.2)^s}.$$

Dividing both sides by  $\log 5$  yields that

$$\left| s \frac{\log 2}{\log 5} - c \right| < \frac{2}{(1.2)^s}.$$

Let  $\gamma := \frac{\log 2}{\log 5}$  and  $[a_0, a_1, a_2, \dots] = [0, 2, 3, 9, 2, \dots]$  be the continued fraction of  $\gamma$ , and let  $p_k/q_k$  be its  $k$ -th convergent. *Mathematica* reveals that

$$q_{23} < 2.5 \cdot 10^{11} < q_{24}.$$

$a_M := \max \{a_i; i = 0, \dots, 24\} = a_{23} = 42$ . By the properties of continued fractions, we obtain

$$\frac{1}{(a_M + 2)s} < \left| s \frac{\log 2}{\log 5} - c \right| < \frac{2}{(1.2)^s}$$

which yields that  $s \leq 45$  as  $a_M = 42$ . Since  $5^c < 2^s$ , then we deduce that  $c \leq 19$ . A quick inspection using *Mathematica* reveals that there is no solution of the equation (4) with  $1 \leq \max \{a, b\} \leq c$  and  $6 \leq c \leq 19$ .  $\square$

### 3 Proof of Theorem 1.1

Firstly, assume that  $1 \leq c \leq 5$ . Then the solution of the equation (2) is given in Theorem 1.1. From now on, suppose that  $c \geq 6$ . Lemma 2.3 gives that  $\alpha^{s(n-2)} < F_n^s < 3 \cdot 5^c < 5^{1.1c}$ . So, we have the fact  $s < c$ . Since  $a, b, c \geq 1$ , then  $n \geq 3$  holds. If  $n = 3$ , then we can rewrite formula (2) as

$$2^s = 2^a + 3^b + 5^c.$$

which is investigated in Lemma 2.4. If  $n = 4$  and  $n = 5$  hold, then we arrive at a contradiction since the left-hand side of the equation  $F_n^s = 2^a + 3^b + 5^c$  is odd, while right-hand side is even. Therefore, we suppose that  $n \geq 6$ .

Using formula (1), we rewrite the equation (2) as

$$F_n^s - 5^c = 2^a + 3^b. \quad (7)$$

Since  $\max \{a, b\} \geq 1$ , then the right-hand side of above equation is positive. Dividing both sides of the equation (7) by  $5^c$ , we obtain

$$\left| F_n^s 5^{-c} - 1 \right| < \frac{2}{5^{0.3c}} \quad (8)$$

where we use  $2 < 3 < 5^{0.7}$ .

In the application of theorem of Matveev, we take  $\alpha_1 = F_n$ ,  $\alpha_2 = 5$ ,  $b_1 = s$ ,  $b_2 = c$ . We also take

$$\Lambda := F_n^s 5^{-c} - 1.$$

Since we assume  $n \geq 6$ , then it is obvious that  $\Lambda \neq 0$ . We can take the degree  $D = 1$ . Then  $A_1 = \log F_n$  and  $A_2 = 1.61 > \log 5$  follow. As  $s < c$ , then we get  $B = c$  together with  $t = 2$ .

After applying the inequality (3) to get lower bound for the form  $\Lambda$ , then we have

$$e^{-C_{2,1}(1+\log c)\times 1.61\times \log F_n} < \frac{2}{5^{0.3c}}, \quad (9)$$

where  $C_{2,1} = 1.4 \times 30^5 \times 2^{4.5}$ . Hence, we obtain that

$$\frac{c}{\log c} < 2.5 \times 10^9 (n-1), \quad (10)$$

where we used the fact  $1 + \log c < 2 \log c$ . It is easy to prove that  $\frac{x}{\log x} < A$  yields  $x < 2A \log A$ . After rewriting the formula (9), we obtain

$$c < 7.3 \times 10^{10} (n-1) \log (n-1) \quad (11)$$

by the inequality  $21.64 + \log (n-1) < 14.6 \times \log (n-1)$ .

Assume that  $n \in [6, 233]$ . Label  $z := s \log F_n - c \log 5$ . Hence, by the equation (8)

$$0 < z < e^z - 1 < \frac{2}{5^{0.3c}} \quad (12)$$

follows. Dividing both sides by  $\log 5$ , we obtain

$$0 < s\gamma - c < 1.25 \times 5^{-0.3c}$$

where  $\gamma := \frac{\log F_n}{\log 5}$ . Let  $[a_0, a_1, a_2, \dots]$  be the continued fraction of  $\gamma$ , and let  $p_k/q_k$  be its  $k$ -th convergent. We have

$$s < c < 9.23 \times 10^{13}$$

by the inequality (11). A quick inspection using *Mathematica* reveals that  $q_{40} > M$ . Moreover,  $a_M := \max \{a_i, i = 0, 1, \dots, 40\} = 3996$ . From the properties of continued fractions, we get that

$$|s\gamma - c| > \frac{1}{(a_M + 2)s}. \quad (13)$$

Comparing the estimates (12) and (13) we get

$$\frac{1}{3998s} < 1.25 \times 5^{-0.3c} \Rightarrow 5^{0.3s} < 5^{0.3c} < 4997.5s,$$

which yields that  $s \leq 23$ . Hence,  $c \leq 1581$  follows. In order to decrease the upper bound for  $c$ , we use that  $\nu_5(F_n^s - 2^a - 3^b) = c$ . Thus, *Mathematica* returns  $\nu_5(F_n^s - 2^a - 3^b) = c \leq 12$  for  $1 \leq s \leq 23$ ,  $6 \leq n \leq 233$  and  $c > \max \{a, b\} \geq 1$ . Therefore,  $c \leq 12$  gives that  $n \leq 44$ . Then the solutions of the equation (2) are given Theorem 1.1.

From now on, assume that  $n > 233$ . In order to find the upper bound for  $c$ , we use the key argument in the paper [8]. Let  $x := \frac{s}{\alpha^{2n}}$ . From the above inequality (11), it follows that

$$x < \frac{7.3 \times 10^{10} (n-1) \log (n-1)}{\alpha^{2n}} < \frac{2}{\alpha^n},$$

where it holds for  $n > 233$ . We now write

$$F_n^s = \frac{\alpha^{ns}}{5^{\frac{s}{2}}} \left( 1 - \frac{(-1)^n}{\alpha^{2n}} \right)^s. \quad (14)$$

In the paper of Luca and Oyono [8], it was proven that

$$\left| \left( 1 - \frac{(-1)^n}{\alpha^{2n}} \right)^s - 1 \right| < \frac{2}{\alpha^n}. \quad (15)$$

Let  $\Lambda_2 := 5^{c+\frac{s}{2}}\alpha^{-ns} - 1$ . From the formulas (2) and (14) together with the inequality (15), we have

$$|\Lambda_2| < \frac{2}{\alpha^n} + \frac{(2^a + 3^b) 5^{\frac{s}{2}}}{\alpha^{ns}}. \quad (16)$$

For the inequality (16) the facts that  $2^a + 3^b < 2 \times 5^{0.7c}$  and  $n > 233$  yield that

$$|\Lambda_2| < 0.8.$$

The last inequality gives that  $\frac{5^{\frac{s}{2}}}{\alpha^{ns}} < \frac{2}{5^c}$ . The inequality (16) yields

$$|\Lambda_2| < \frac{2}{\alpha^c} + \frac{2}{\alpha^n} = \frac{4}{\alpha^l},$$

where  $l = \min\{n, c\}$ . We use again the theorem of Matveev. We take  $k = 2$ ,  $\alpha_1 := \alpha$ ,  $\alpha_2 := 5$ ,  $b_1 := ns$ ,  $b_2 := c + \frac{s}{2}$ . As in the previous application of Matveev's result, we can take  $D := 2$ ,  $A_1 := 0.5$ ,  $A_2 := 1.61$ . Note that  $\alpha^c < 5^c < \alpha^{s(n-1)}$  gives  $c + \frac{s}{2} < c + s < ns$ . So, we take  $B := ns$ . We thus get that

$$\exp(-C_{2,2}(1 + \log ns) \times 0.5 \times 1.61) < \frac{4}{\alpha^l},$$

where  $C_{2,2} = 1.4 \times 30^5 \times 2^{4.5} \times 4(1 + \log 2)$ . This leads to

$$l < \frac{C_{2,2}(\log ns) \times 1.61}{\log \alpha}.$$

If  $l = n$ , then the last inequalities

$$\begin{aligned} n &< \frac{C_{2,2}(\log ns) \times 1.61}{\log \alpha} \\ &< \frac{C_{2,2}(\log n (7.3 \times 10^{10} (n-1) \log(n-1))) \times 1.61}{\log \alpha} \\ &< \frac{C_{2,2}(\log(7 \times 10^{10} n^3)) \times 1.61}{\log \alpha} \end{aligned}$$

give that  $n < 1.92 \times 10^{12}$ . By the inequality (11), we get

$$c < 7.3 \times 10^{10} (n-1) \log(n-1) < 4 \times 10^{24}.$$

If  $l = c$ , then we have that

$$c < \frac{C_{2,2}(\log ns) \times 1.61}{\log \alpha} < \frac{C_{2,2}(\log 6.8c) \times 1.61}{\log \alpha}$$

yields  $c < 6 \times 10^{11}$ , where we used the fact

$$\alpha^{ms} < 5^{1.2c} \alpha^{2s} < \alpha^{4.8c} \alpha^{2s} = \alpha^{6.8c}.$$

At any rate, we get

$$c < 4 \times 10^{24}.$$

Next we take  $\Gamma := \left(\frac{s}{2} + c\right) \log 5 - ns \log \alpha$ . Observe that  $\Lambda_2 = e^\Gamma - 1$ . Since  $|\Lambda_2| < 0.8$ , then we have  $|e^\Gamma - 1| < 0.8$ , which yields that  $e^{|\Gamma|} < 2$ . Hence,

$$|\Gamma| \leq e^{|\Gamma|} |e^\Gamma - 1| < 2 |\Lambda_2| < \frac{2}{\alpha^c} + \frac{2}{\alpha^n}.$$

This leads to

$$\begin{aligned} \left| \frac{\log \alpha}{\log 5} - \frac{c + \frac{s}{2}}{ns} \right| &= \left| \frac{\log \alpha}{\log 5} - \frac{2c + s}{2ns} \right| < \frac{1}{ns \log 5} \left( \frac{2}{\alpha^c} + \frac{2}{\alpha^n} \right) \\ &< \frac{1}{374s} \left( \frac{2}{\alpha^c} + \frac{2}{\alpha^n} \right) \end{aligned} \quad (17)$$

since  $n > 233$ . Assume that  $c \geq 30$ . In this case, note that  $\alpha^n > 160c^2$  (as  $n < ns < 6.8c$ ) and  $\alpha^c > 160c^2$ . Hence, we get that by the inequality (17) by the fact  $\alpha^{ns} < \alpha^{6.8c}$ ,

$$\begin{aligned} \left| \frac{\log \alpha}{\log 5} - \frac{2c + s}{2ns} \right| &< \frac{1}{14960sc^2} < \frac{1}{14960c^2} \\ &< \frac{6.8^2}{14960(ns)^2} < \frac{1}{80(2ns)^2}. \end{aligned} \quad (18)$$

By a criterion of Legendre, the rational number  $\frac{2c + s}{2ns}$  converts to  $\gamma := \frac{\log \alpha}{\log 5}$ .

Let  $[a_0, a_1, a_2, \dots] = [0, 3, 2, 1, \dots]$  be the continued fraction of  $\gamma$ , and  $p_k/q_k$  be its  $k$ -th convergent. Assume that  $\frac{2c + s}{2ns} = \frac{p_t}{q_t}$  for some  $t$ . We have  $q_{49} > 4 \times 10^{24}$ . Thus,  $t \in \{0, 1, \dots, 49\}$ . Furthermore,  $a_k \leq 59$ , for  $k = 0, 1, \dots, 49$ . From the well-known properties of continued fractions, we get that

$$\left| \gamma - \frac{2c + s}{2ns} \right| = \left| \gamma - \frac{p_t}{q_t} \right| > \frac{1}{(a_t + 2)q_t^2} > \frac{1}{(a_t + 2)(2ns)^2} > \frac{1}{4 \times 61 \times 6.8^2 \times c^2}. \quad (19)$$

After combining the inequalities (18) and (19), then

$$\frac{1}{4 \times 61 \times 6.8^2 \times c^2} < \frac{1}{14960c^2s}$$

gives  $s < 1$ . But, this is not possible.

Therefore,  $c$  is at most 29. By the inequality  $ns < 6.8c$ , we obtain that  $ns \leq 197$ , which is false as  $n > 233$ .

Hence, the proof theorem is completed.

## Acknowledgements

The authors are grateful to the reviewers for their helpful and valuable comments and remarks.

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